Graph-IQ⟩⟨Cl, a Graph-Based Quantum/Classical Algorithm for Efficient Electronic Structure on Hybrid Quantum/Classical Hardware Systems: Improved Quantum Circuit Depth Performance

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ABSTRACT: We present a procedure to reduce the depth of quantum circuits and improve the accuracy of results in computing post-Hartree–Fock electronic structure energies in large molecular systems. The method is based on molecular fragmentation where a molecular system is divided into overlapping fragments through a graph-theoretic procedure. This allows us to create a set of projection operators that decompose the unitary evolution of the full system into separate operators that can be asynchronously spawned onto a potentially large ensemble of classical and quantum hardware systems. We demonstrate this method by computing Unitary Coupled Cluster Singles and Doubles (UCCSD) energies for a set of [H₂]ₙ clusters, with n ranging from 4 to 128. We implement our methodology using quantum circuits, and when these quantum circuits are processed on a quantum simulator, we obtain energies in agreement with the UCCSD energies in the milli-hartree energy range. We also show that our circuit decomposition approach yields up to 9 orders of magnitude reduction in the number of CNOT gates and quantum circuit depth for the large-sized clusters when compared to a standard quantum circuit implementation available on IBM’s Quantum Information Science kit, known as Qiskit.

1. INTRODUCTION

Computing accurate electronic properties, including the effect of electron correlation and nuclear dynamics, is at the heart of modern quantum chemistry, with potential impact on materials discovery,¹−³ and the accurate study of biological⁴−¹⁰ and atmospheric¹¹−¹³ processes. However, such studies are deeply confounded by (a) the steep (algebraic) computational scaling of accurate electron correlation methods,⁵−⁷ where, for example, the gold standard of electronic structure theory, namely, CCSD(T), scales as O(N⁶−⁷) and (b) the likely exponential scaling of quantum nuclear dynamics. Over the years, several classical algorithms have been developed to improve the computational scaling of both problems.¹⁷−²¹ For example, molecular fragmentation has recently become a critical tool to compute electronic properties and has grown to provide extremely accurate and effective computational paradigms. Similarly, the effect of nuclear dynamics has also been an active area of study for complex reactive processes.²²−²⁴

Recently, multiple quantum computing technologies, such as ion traps,²⁶−²⁹ superconducting coils,³⁰−³¹ Bosonic processors with photons,³²−³³ solid-state devices and quantum dots inside cavities,³⁵−³⁸ and Rydberg atoms,³⁹−⁴¹ have emerged as potential alternative computational platforms to address complex computational challenges. Additionally, algorithms to approximate electron correlation problems,⁴²−⁶⁵ for small molecular systems, and quantum nuclear dynamics problems⁴⁴−⁷¹ have been implemented on quantum hardware devices using trapped atomic ions, photons, nuclear spins, quantum dots, Rydberg atoms, and superconducting circuits. However, as outlined in Figure 1, application of standard quantum circuit models²⁵ to treat electronic structure problems leads to a rapid increase in the circuit depth and the number of CNOT gates. This contributes greatly to the accumulated error during quantum propagation. Quantum gate fidelity for CNOT gates is generally of the order of 95%,⁷² and the improvement of such gate fidelity is an active area of quantum hardware development.⁷³ This is related to the fact that the CNOT gates require maximally entangling XX(π/2) gates and hence have lower-quality performance⁷² as compared to small-angle XX gates, and this aspect contributes to error propagation. As seen in Figure 1, the number of quantum gates, the number of CNOT gates, and the circuit depth increase exponentially with system size. This leads to a dramatic increase in the error in quantum
propagation thus restricting both the size and quality (in terms of basis set sizes that can be routinely used) of performance. Thus, despite the growing set of available quantum hardware platforms and the accompanying set of complex quantum algorithms, performing accurate, state-of-the-art quantum chemical calculations will remain a significant challenge for the foreseeable future. In this publication, we present a new hybrid quantum/classical algorithm that helps reduce the number of CNOT gates, the circuit depth, and the total number of gates by several orders of magnitude. This paper outlines a new hybrid algorithm based on a graph-theoretic approach to molecular fragmentation and is geared toward performing electron correlation calculations, potentially on an ensemble of quantum and classical hardware systems. From here on, the algorithm studied here is referred to as the “Graph-IQ” (GI) algorithm since it contains an independent set of classical and quantum algorithmic components inside a single umbrella. That is, the overall computational workload is partitioned, through graph theory based on computational complexity analysis, into (a) classical computing sections that are carried out on traditional classical electronic structure packages, such as Gaussian,74 Psi4,75 ORCA,76 and Quantum ESPRESSO,77 and (b) quantum computing sections that are carried out using quantum circuit models. Furthermore, the Graph-IQ (GI) algorithm is quantum hardware-agnostic and is developed with the goal to be implemented on all quantum hardware technologies, and, in fact, is designed to be used on an ensemble of such quantum hardware systems for any given calculation. In essence, our Graph-IQ (GI) algorithm yields a new approach that reduces the required quantum circuit depth, the number of quantum gates, and the number of CNOT gates (by several orders of magnitude) that contribute to error accumulation (the scaling of these three properties is shown in Figure 1), through a graph-theory-based projection operator formalism. Thus, given this reduction, our algorithm, which is complementary to other ideas in the literature22,23,81,82 that attempt to reduce quantum gate complexity, potentially improves the quantum algorithmic efficiency, provides a new avenue for quantum resource management, and also reduces the accumulation of errors during the demonstrated electronic structure calculations on quantum hardware. Given the limitations of quantum circuit gate fidelities within the gate model, this algorithm, which we expect, will become a central piece in the quantum/classical computing of chemical systems.

Our approach begins with a graph-theoretic molecular fragmentation procedure introduced in refs 22, 23, and 78–85, where the key idea begins with the well-known ONIOM methodology,86 however, then, the “model” and “real” system energies and gradients within ONIOM are now constructed using many-body expansions up to an arbitrary rank. Furthermore, these many-body expansions are obtained in a general fashion using adaptive graph-theoretic techniques that are computationally available within standard protocols in the Python programming language. This paper is organized as follows: in Sections 2 and 3, we outline our graph-theoretic molecular fragmentation approach and also describe how these are to be implemented within a hybrid ensemble of quantum and classical hardware systems to yield reduced complexity of quantum circuits. We show that the graph-theoretic approach yields a unitary transformation applied to the quantum circuit model corresponding to the full quantum molecular system to reduce it into a family of decoupled, parallel quantum circuit models, each of which has much lower complexity as compared to the parent circuit. Thus, we expect that the error propagation in this new family of circuits is far less as compared to that belonging to the parent full system. In Section 4, we demonstrate the reduction in quantum circuit complexity scaling that arises from our algorithm through results for hydrogen molecular clusters. Conclusions are given in Section 5.

2. GRAPHICAL REPRESENTATION OF LOCAL MANY-BODY INTERACTIONS MAPPED TO QUANTUM CIRCUIT MODELS

Thus far, the electronic structure of small systems42–46,50,87 or low-classical scaling approximations of the same for larger systems22 have been implemented on quantum technologies. Here, we develop an approach, based on graph-theoretic fragmentation of molecular systems, that provides a pathway toward an accurate depiction of electron correlation on large systems using hybrid quantum/classical algorithms appropriate for an ensemble of quantum and classical computing environments. In a series of publications,22,23,78–85 we have developed graph theory-based techniques to compute efficient and adaptive many-body expansions that are then embedded within the ONIOM86 multi-layer approach; the accuracy and efficiency of these methods have been demonstrated on classical hardware systems. In this paper, we develop algorithms that make this approach applicable to a hybrid set of quantum and classical hardware systems. The salient features of this approach are as follows: the molecular assembly is partitioned into a set of nodes or vertices. Generally, these nodes are determined on a chemical basis and may include groups such as water and hydro-}

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**Figure 1.** Quantum circuit depth complexity as a function of system size is illustrated here for a family of $(H_2)_n$ clusters. As the system size grows, a standard quantum circuit implementation afforded by the commonly used Qiskit quantum computing model becomes prohibitively complex.
that is now written in a general and computationally robust lower level of theory, where the perturbative correction is the Consistent with the notions behind ONIOM, we energetic measure we begin with 22,23,78

Together, these nodes and edges de

graph. This graph presents a distance-based truncation of expensive where the left side, \( \Omega_{\text{MBE}}(\mathbf{x}) \), denotes the graph-theoretically obtained many-body correction to ONIOM, and the term \( E_{\text{MBE}}^{\text{level,1}}(\mathbf{x}) \) on the right side may encompass the full system or a chosen “active site”, and we have considered both options within AIMD78–82,85 and quantum nuclear potential surface treatments.22,23 In addition to the extrapolatory, ONIOM-like form of eq 2, each term in the extrapolation is a many-body expansion, that now written in a general and computationally robust fashion up to order (or rank) \( R \) as

\[
E_{\text{MBE}}^{\text{level,1}} = \sum_{r=0}^{R} \sum_{\alpha} (-1)^{\alpha} \sum_{m=\alpha}^{R} (-1)^{m} p_{m}^{r,m,\alpha}
\]

where \( p_{m}^{r,m,\alpha} \) is the number of times the \( \alpha \)th \((r+1)\)-body term appears in all \((m+1)\)-body terms (in set \( \mathbf{V}_{m} \) for \( m \geq r \), and consequently, \( \Omega_{\text{MBE}} \equiv \left[ \sum_{m=r}^{R} (-1)^{m} p_{m}^{r,m,\alpha} \right] \) is the overcounting correction for the number of times the \( \alpha \)th \((r+1)\)-body term appears in all objects of rank greater than or equal to \( r \). It is important to emphasize that eq 3 is essentially identical to standard many-body expressions but presented now using graph theory. Thus, the full energy expression, which combines eqs 2 and 3, becomes

\[
E_{\text{MBE}}^{\text{ONIOM}}(\mathbf{x}) = E_{\text{MBE}}^{\text{level,0}} + \sum_{r=0}^{R} (-1)^{r} \sum_{\alpha} \left( \sum_{m=\alpha}^{R} E_{\text{MBE}}^{\text{level,1}}(\mathbf{V}_{\alpha}) - E_{\text{MBE}}^{\text{level,0}}(\mathbf{V}_{\alpha}) \right) \Omega_{\text{MBE}}(\mathbf{x})
\]

In practice, the individual simplexes in eqs 3 and 4 are computed independently and asynchronously, leading to a general Python-based parallel implementation. Furthermore, our implementation also allows use of separate electronic structure packages for each level of treatment and currently supports the simultaneous use of Gaussian,71 ORCA,76 Psi4,75 Quantum ESPRESSO,77 and OpenMX.88 A few key ideas regarding our graph-theoretic formalism are presented in the Supporting Information to facilitate the discussion in the paper.

3. GRAPH-THEORETICALLY DETERMINED PARALLEL QUANTUM CIRCUITS FROM EQS 3 AND 4

As seen in the Introduction section, the problem of quantum circuit depth is a serious limitation in obtaining accurate quantum computing results. In this section, we present a general solution to this problem based on the graph-theoretic paradigm presented in the previous section. The graphical decomposition of the instantaneous molecular structure is parsed, and the steep scaling aspects (represented as “level, 1” in eq 4) are spawned to a quantum computing or quantum simulation system, while the lower scaling components are retained on classical hardware. To achieve this, in Section 3.1, we develop a general projection operator-based formalism, which when applied to a full-system molecular quantum circuit, as shown in in Section 3.2, reduces the circuit into a family of parallel quantum circuits. In the Results and Discussion section, we have then shown that this approach reduces the quantum circuit depth problem, shown in Figure 1, by several orders of magnitude.

3.1. Graph-Theory-Based Projection Operators. In this section, we introduce a new quantum circuit decomposition technique that reduces the complexity of circuits (measured in terms of circuit depth and number of CNOT gates as seen in Figure 1) by several orders of magnitude. We begin with a Hilbert space decomposition scheme using the set-theoretic inclusion–exclusion principle.89 The projection technique is then adapted to a graph problem, which in the next subsection is used to decompose any arbitrary unitary operator (or quantum circuit) into parallel, but overlapping, streams of computing processes that can be executed on a cluster of quantum and classical hardware systems.

Let us begin with a Venn diagram that divides a coordinate representation \( \mathbf{x} \) into regions, \( A, B, C, \) and so forth. The regions may intersect, and in Figure 3a, for concreteness, we have superimposed the Venn diagram on top of a \( \text{H}_2 \) molecular cluster. Thus, the Venn divides a molecular system into several regions. Using the principle of inclusion and exclusion,89 the resolution of the identity for the Hilbert space depicted within the Venn diagram may be written as

\[
\mathbf{x} = A \cup B \cup C \]

\[
\mathbf{x} = \sum_{r=0}^{R} \sum_{\alpha} (-1)^{\alpha} \sum_{m=\alpha}^{R} E_{\text{MBE}}^{\text{level,1}}(\mathbf{V}_{\alpha}) - E_{\text{MBE}}^{\text{level,0}}(\mathbf{V}_{\alpha}) \Omega_{\text{MBE}}(\mathbf{x})
\]

where \( p_{m}^{r,m,\alpha} \) is the number of times the \( \alpha \)th \((r+1)\)-body term (in set \( \mathbf{V}_{\alpha} \)) appears in all \((m+1)\)-body terms (in set \( \mathbf{V}_{m} \) for \( m \geq r \)).
where the left side is the identity since it sums over the entire Hilbert space represented here by sets $A$, $B$, and $C$, and integrals involving the dyadic terms, $|x⟩⟨x|$, are within a chosen set depicted within the Venn diagram. Additionally, we have also introduced projection operators

$$\mathcal{P}_A \equiv \int_A dx \, |x⟩⟨x|$$

(6)

that project out portions of a Hilbert space depicted in eq 5 and only include the portion of the diadic sum within a given set. While eq 5 arises from the principle of inclusion and exclusion well known in set theory and can be generalized to

Figure 3. In figure (b), we provide an illustration of the sets $A$, $B$, and $C$ for the system in figure (a). Figure (b) is used to construct eq 5. In figure (c), the sets are provided a graph-theoretic form to make the transition from eq 5 to 7 clear. Figure (c) can be realized through eq 9 which is a special case of eq 7.

Figure 4. Quantum circuit decomposition based on graphical partitioning of molecular systems. The second-quantized Hamiltonians for node/edge/face fragments, with molecular orbitals obtained from Hartree–Fock, are used as the input for the Qiskit quantum computing framework45 to obtain quantum circuit models shown on the right side of the figure, one circuit for each fragment. Details of this process are given in the Results and Discussion section.
simplexes as of projectors that encompass nodes, edges, and higher-order (also see Figure 2). As usual, the graph comprises vertices, edges, and rank- \( r \) simplexes. The graph decomposition of molecular systems. This helps compute \( \{ E^{\text{level}, l} \} \) in eq 4. To compute the lower-scaling \( \{ E^{\text{level}, l} \} \), we use classical computing algorithms as seen on the left side of the figure here.

an arbitrary number of sets, an alternate approach to divide the space represented by the identity operator, \( I \), can be obtained by reintroducing the graph \( G \equiv \{ V_0, V_1 \} \) from the previous section (also see Figure 2). As usual, the graph comprises vertices, \( V_0 \) edges, \( V_1 \), and rank- \( r \) simplexes. An equivalent expression for the resolution of the identity in eq 5 may now be obtained in terms of projectors that encompass nodes, edges, and higher-order simplexes as

\[
I = \sum_{a \in V_0} M_{a,0}^0 \mathcal{P}_a - \sum_{a \in V_1} M_{a,1}^1 \mathcal{P}_a + \sum_{a \in V_2} M_{a,2}^2 \mathcal{P}_a - \cdots
\]

\[
= \sum_{r=0}^{R} (-1)^r \sum_{a \in V_r} M_{a,r}^r \mathcal{P}_a
\]

(7)

Here, as in eqs 3 and 4, \( M_{a,r}^r \equiv \sum_{m=r}^{R} (-1)^m \mathcal{P}_a^{r,m} \) and \( \mathcal{P}_a^{r,m} \) are the number of times the \( \alpha \)-th rank- \( r \) term (in set \( V_r \)) appears in all rank- \( m \) terms (in set \( V_m \)), for \( m \geq r \). Consequently, \( M_{a,r}^r \) is the overcounting correction for the number of times the \( \alpha \)-th rank- \( r \) term appears in all objects of a rank greater than or equal to \( r \). It must be noted that \( \mathcal{P}_a^{r,m} \) are also the number of supersets of the \( \alpha \)-th rank- \( r \) simplex, and the projectors, \( \mathcal{P}_a^{r,m} \), yield the \( \alpha \)-th rank- \( r \) simplex. These projection operators are used in the next section to reduce the quantum circuit depth.

The parallels between eqs 7 and 5 may be further explored by rewriting eq 7 in a decreasing order of rank, that is

\[
I = (-1)^R \sum_{a \in V_R} \mathcal{P}_a - \sum_{a \in V_{R-1}} M_{a,R-1}^{R-1} \mathcal{P}_a - \sum_{a \in V_{R-2}} M_{a,R-2}^{R-2} \mathcal{P}_a - \cdots
\]

(8)

where the appearance of alternating signs is clear and resembles that in eq 5, and the factor \( M_{a,R-1}^{R-1} \) is the number of times the \( \alpha \)-th rank- \( \{ R - 1 \} \) simplex appears in all rank- \( R \) simplexes. Additionally, for \( R = 1 \), eq 8 becomes

\[
I = - \sum_{a \in V_1} \mathcal{P}_a + \sum_{a \in V_0} M_{a,0}^0 \mathcal{P}_a
\]

(9)

which, for the graph in Figure 3c, leads to an identical result as in eq 5, constructed for Figure 3b.

3.2. Quantum Circuit Depth Reduction through Parallel Quantum Processing Using the Projection Operators, \( \{ \mathcal{P}_a \} \). We now begin with a quantum circuit depicted using the symbol \( \mathcal{U} \), a unitary operator, that propagates a state that represents the electronic structure for the full molecular system on some quantum hardware system. However, as noted in Figure 1, the complexity of such a circuit and the associated resources may grow rapidly as system size grows. To overcome this issue, we may apply the graph-theoretically defined resolution of identity in eq 7 to \( \mathcal{U} \) to decompose it into a family of parallel quantum circuits given by

\[
2889 \quad \text{https://doi.org/10.1021/acs.jctc.1c01303}
\]

It is critical to emphasize at this stage that it is now accepted\textsuperscript{90,91} that universal, fully fault-tolerant quantum computers are a rather distant dream\textsuperscript{90} and new frontiers such as Noisy Intermediate-Scale Quantum (NISQ)\textsuperscript{90} systems have emerged. To effectively use such NISQ machines with their limited coherence times, a hybrid approach that interleaves NISQ machines with classical computers has been proposed in ref \textsuperscript{91}. This is complemented by orthogonal developments in quantum chemical algorithm developments where the variational quantum eigensolver (VQE)\textsuperscript{62,63} is already known to be a hybrid quantum/classical approach. In this regard, our approach here, containing quantum and classical algorithmic components, is a contribution that can be used for ground-state quantum chemical calculations at enhanced accuracy with lower computational complexity\textsuperscript{80\textendash}83,85 on a hybrid stream of quantum and classical systems. The algorithm is built to be asynchronous and parallel, and these computational aspects will be considered in a future publication.

In general, a molecular system, and in future a molecular surface,\textsuperscript{83} is coarse-grained to create a graphical representation as explained at the top of Section 2. This yields simplexes that are then used to create a “bucket of fragments” as referred to in Figures 4 and 5. Each fragment needs to be used to compute energies at two levels of theory, \( E_{\text{ UD}} \) and \( E_{\text{ UCCSD}} \). As stated in Figure 5, the quantity \( E_{\text{ UD}} \) is computed on classical hardware systems, whereas the quantity \( E_{\text{ UCCSD}} \) is to be computed on quantum hardware using the quantum circuits \( \mathcal{U}_{\alpha,r} \) as facilitated by the graph-theoretic partitioning method, thus expanding the realm of applications for quantum computation. In this publication, we do not present results that use actual quantum hardware, but we present results where the family of quantum circuits, \( \{ \mathcal{U}_{\alpha,r} \} \), is processed using the Qiskit system\textsuperscript{25} but implemented on a quantum simulator.\textsuperscript{25} Further computational details are presented in Section 4.1.

4. RESULTS AND DISCUSSION

In order to gauge the accuracy and reduction in computational complexity arising from our Graph-\( \{Q\}\) method presented above, we have applied this approach to a range of hydrogen molecular cluster problems. These systems are critical for applications related to energy storage.\textsuperscript{92\textendash}97 In particular, the safe and efficient storage\textsuperscript{98,99} of molecular hydrogen is of paramount importance to potential developments in new fuel cell technologies.\textsuperscript{101\textendash}109 Furthermore, the study of ortho- and para-hydrogen\textsuperscript{104\textendash}109 at low-temperatures has been a funda-

\[ \mathcal{U} = \sum_{r=0}^{R} (-1)^{r} \sum_{\alpha \in \mathcal{V}^{r}} \mathcal{M}_{\alpha}^{r} \{ \mathcal{P}_{\alpha,r}, \mathcal{U} \} \]

where

\[ \{ \mathcal{U}_{\alpha,r} = \mathcal{P}_{\alpha,r}, \mathcal{U} \} \]

represents here a set of projected quantum circuits, one for each simplex. When a molecular system is divided using the graph, as presented in Section 2, the set \( \{ \mathcal{U}_{\alpha,r} \} \) yields one quantum circuit for each molecular fragment. This is illustrated in Figure 4. Thus, using eq 10, it is possible to independently construct the unitary quantum circuits, \( \mathcal{U}_{\alpha,r} \), one for each molecular fragment and use these in parallel on a family of decoupled quantum hardware systems to perform the computations needed to obtain the approximations in eqs 3 and 4.

In this publication, we use the individual quantum circuits, \( \{ \mathcal{U}_{\alpha,r} \} \), to obtain the family of fragment energies \( \{ E_{\text{ UD}}^{\text{ frag}} \} \), using quantum circuit models, where level, 1 here is Unitary Coupled Cluster Singles and Doubles (UCCSD), which when used in eq 10, or equivalently, eq 3, yields an approach to compute the left side of eq 4, using a stream of parallel processes. This idea is presented in Figure 4. Because these independent circuits, \( \{ \mathcal{U}_{\alpha,r} \} \), are for much smaller fragments as compared to the full system of interest, one may find the error propagation to be limited. We indeed find this to be the case for the examples discussed later.

Furthermore, eq 4 contains both level, 1 calculations and level, 0 calculations. This kind of composite approach has been shown to converge faster as a function of maximum rank \( R \), for ground-state post-Hartree–Fock energies, AIMD trajectories, and multi-dimensional potential calculations using post-Hartree–Fock energies and gradients. Here, level, 1 is UCCSD, whereas level, 0 will include a computationally less expensive DFT approximations. Thus, the algorithm presented here envisions spawning out a family of processes on an ensemble of classical and quantum systems, and this process is depicted in Figure 5 and, in more detail, in Figure 6.
while the key properties of such systems involve the detailed study of electronic as well as nuclear degrees of freedom, in this publication, our goal is to gauge the accuracy of the approach introduced above the ability to reduce the complexity of quantum circuits in obtaining post-Hartree–Fock (coupled cluster)-level electronic energies for such systems. Thus, we compute post-Hartree–Fock electronic energies using the quantum circuit decomposition algorithm mentioned above and compare these results with those obtained using classical algorithms obtained at the unitary coupled cluster level of theory (UCCSD).

Hydrogen molecular clusters of various sizes, \((\text{H}_2)_\text{a} (\text{a} = 2, ..., 128)\) have been treated with the Graph-QCI method to demonstrate scalability and accuracy as the system size grows. Specifically, the analysis of errors due to the truncation in the rank \(R\) in eq 10 and edge length cutoff used in the graph definition allows us to gauge the fragment circuit complexity needed to achieve an acceptable (milli-hartree) level of accuracy.

4.1. Computational Aspects. Figure 5 provides a brief overview of the algorithm which is further elaborated upon in Figure 6. Specifically, simplexes labeled \((\alpha, r)\), in eqs 10 and 4, are pre-determined on classical hardware. Each hydrogen molecule is treated as a node in the graph. The process of defining a graph is controlled by two parameters, the maximum edge length cutoff and the maximum order (or rank) of the many-body terms which is represented in eqs 10 and 4 using the symbol \(R\). As the maximum edge length increases, the fragments (including higher-order simplexes) grow in number and size rapidly, thus increasing computational complexity while also presenting a significant challenge for parallel processing. The edge length cutoff may be chosen to be high enough to include all critical interactions. The maximum edge length has been chosen here based on previous studies, as well as additional analysis presented in the next subsection, and leads to the type of graph presented in Figures 4 and 5.

Once the graph is generated, the system is decomposed into a set of simplexes, or fragments, that essentially now can be processed in a completely independent way. A stream of simplexes is then generated with the goal to compute \(\{E_{\text{level}}(\alpha, r)\}\), \(E_{\text{level}}(\alpha, r)\), and \(E_{\text{level}}\) for the full system. As discussed in ref 85, the stream of fragments generated can be processed in an asynchronous (non-blocking) and parallel manner, to generate the fragment energies \(E_{\text{level}}\) and \(E_{\text{level}}\), to be used in eq 4. In refs 83 and 85, this is done through an MPI parallelized hybrid C++/Python module, which is capable of using multiple electronic structure packages within a single AIMD step and isolated electronic structure calculations conducted during potential surface evaluations.

The implementation of the approach in classical computing platforms currently supports the following set of external electronic structure packages during a single energy and gradient evaluation: Gaussian, ORCA, and Ps4, for molecular and cluster calculations, and Quantum ESPRESSO and OpenMX for condensed phase studies. Here, this aspect is further expanded to also include the use of quantum algorithms to determine \(E_{\text{level}}\). (See Figures 5 and 6.) All level, 0 calculations shown in eq 4 were performed at the UCCSD level of theory, on a quantum simulator using an appropriate quantum circuit model. This choice of complex-levels of theory has been demonstrated to be accurate and to yield a rapidly convergent sequence of results with increasing graphical rank \(R\). The choice of basis sets for all calculations in this publication is limited to the minimal STO-3G since a larger basis set introduces more Hartree–Fock orbitals, which are mapped to an even greater number of excitation operators, thus significantly increasing the number of qubits required as well as the associated quantum circuit depth. While our approach will address all of these challenges, by reducing the needed circuit depth, in this publication, we probe the principle behind our execution model for accuracy purposes. Furthermore, as we will see later, it is straightforward to integrate the current scheme into other circuit optimization techniques, to further increase the size of systems that can be studied, and this will be considered as a part of future publications.

The following steps are used to obtain \(\{E_{\text{level}}(\alpha, r)\}\) with quantum algorithms:

1. Hartree–Fock orbitals for each molecular fragment are classically pre-computed to create a family of second-quantized Fermionic Hamiltonians, and these are together provided as input to the Qiskit quantum computing framework, as shown on the bottom-right portion of Figure 6.

2. The family of fragment second-quantized Fermionic Hamiltonians is then converted into quantum circuits, \(\{U_{\alpha, r}\}\), using a parity mapping protocol.

3. Each fragment quantum circuit is then executed on the Qiskit’s built-in state vector simulator without using a noise model, and the resultant energies are optimized using the SLSQP optimizer available within the VQE. This yields the family of UCCSD energies, \(\{E_{\text{level}}(\alpha, r)\}\), used to obtain the molecular energy as per eq 4.

The following comments are in order with respect to the steps mentioned above. The hybrid quantum/classical formalism developed here is transparent to the underlying quantum mapping protocol used to convert the Fermionic Hamiltonians for each individual molecular fragment into quantum circuits. Thus, there are no restrictions on the kind of map used to treat each individual fragment Hamiltonian, and in principle, we can use the Jordan–Wigner, Bravyi–Kitaev, or parity mapping transformations. Here, we have tested the Jordan–Wigner scheme as well as the parity mapping protocol, and we found that the latter does in fact reduce the qubit resources needed as suggested in ref 114. Specifically, for each quantum circuit, two qubits are reduced with the spin-parity symmetries of the system, as a separate optimization technique.

There have been several recent studies that allow the implementation of variational quantum algorithms such as the VQE, in combination with an appropriate trial wavefunction such as the UCC ansatz on quantum simulation systems. However, our approach is general enough to allow the possibility to include other eigensolvers apart from the VQE, such as the recently developed contracted quantum eigensolver to compute the set of \(\{E_{\text{level}}(\alpha, r)\}\) values from the associated two-particle reduced density matrices. These aspects will be further investigated in future publications.

4.2. Reduction of Circuit Depth and Associated Accuracy. The complexity of our algorithm is determined by the maximum value of \(R\) within the family of quantum circuits: \(\{U_{\alpha, r}\}_{\alpha, r=0, ..., R}\). Furthermore, these graph-based circuit complex-
ities are also dictated by the maximum edge length that is used to create the graphs. Thus, we first conducted a detailed analysis of the accuracy of our algorithm when both the maximum edge length as well as the rank are varied and chosen to have values up to their respective maximum possible values. We perform this analysis for systems in the size range \( n = 1, \ldots, 8 \), and present our results in Figure 7. For comparison, we provide results where quantum circuits are used for \( \alpha = \) within the algorithm described above and depicted in Figure 6 and when classical algorithms are used to compute the \( \{E^{\text{level},1}(\alpha, r)\} \) energies. For all these cases, the maximum edge length is chosen such that all nodes are connected, that is, completely connected graphs are used, and the rank is progressively increased to study the accuracy as well as the efficiency in Figure 7. For all cases, given the similarity of the behavior of the three measures depicted in Figure 1, namely, the number of quantum gates, the number of CNOT gates, and the circuit depth, and given the sensitivity of CNOT gate fidelity as discussed in ref 46, for the remaining portion of this paper, we use the number of CNOT gates within a quantum circuit as a measure of computational complexity and extent of error propagation.

As noted in Figure 7a, for each system shown along the horizontal axis, the accuracy increases as the maximum rank is increased and is already within 1/100th of a milli-hartree from the correct result at \( R = 2 \). In fact, for all cases in Figure 7a, \( R = 1 \) provides results with acceptable accuracy. In Figure 7b, we present the computational effort needed to obtain the results in Figure 7a. That is, the larger rank calculations need an exponentially greater number of CNOT gates, which is a measurement of computational complexity, as seen in Figure 1. However, Figure 7b should be compared with Figure 7c, where the full-system computational complexity (or CNOT gate count) is presented. The number of CNOT gates in 7c is several orders of magnitude greater as compared to the truncated \( R \).
-level calculations in Figure 7b; for example, using only $\mathcal{R} = 1$ results in a 1–3 orders of magnitude reduction in the number of CNOT gates, as seen by comparison of Figure 7b,c, while maintaining accuracy to within 0.15 milli-hartree. Furthermore, the complexity for a specific choice of rank $\mathcal{R}$, which is dictated by an expected accuracy for the result as shown in Figure 7a, is constant across systems of all sizes. For example, for [H$_2$]$_n$ at $\mathcal{R} = 1$, instead of computing all interactions with the full system circuit, only 1-body and 2-body interactions are accounted for by a set of circuits representing [H$_2$] and [H$_2$]$_2$. The approximation takes advantage of the local nature of chemical systems.

However, there are some caveats. The computational expense of an individual fragment circuit increases with increasing values of $\mathcal{R}$, as seen in Figure 7b. In addition, as shown in Figure 8 and illustrated in Figure 9 for [H$_2$]$_{128}$, the number of fragment circuits also drastically escalates with respect to $\mathcal{R}$. Thus, the benefit of complexity (or the CNOT count) reduction from using fragments will also require the stream of quantum and classical computations to be executed in parallel to attain computational efficiency.

In this publication, we aim to achieve a milli-hartree-level agreement between full UCCSD calculations and those obtained from the algorithms presented here. Thus, based on our results in Figure 7a, we now investigate larger [H$_2$]$_n$ clusters, with $n = 8$, ..., 128, using $\mathcal{R} = 1$ as the maximum rank. The maximum edge length for these larger clusters was determined based on the distribution of all edge lengths presented in Figure 10. This figure indicates that a maximum edge length of 4.0 Å includes most of the short-range interactions. Additional longer-range interactions may also be included by raising the maximum edge length to 7.5 Å; however, as stated above, the number of simplexes grows rapidly as the edge length increases thus raising the complexity of the algorithm.

The set of results for larger H$_2$ clusters is presented in Figures 11 and 12. While Figure 11a summarizes the results for [H$_2$]$_n$ clusters, with $n = 1$, ..., 8, Figure 11b extends this out to [H$_2$]$_{64}$. As the size of the molecular system grows, the corresponding
circuit complexity, as represented by the number of CNOT gates (red dots), increases exponentially. However, the complexity of the current algorithm, as represented by the yellow dots in Figure 11, grows in a gradual fashion while maintaining accuracy in the milli-hartree range as can be seen from the right vertical axis of Figure 11. Clearly, these figures also show that while $g = 1$ appears sufficient and does provide accurate results, the error grows with system size. Thus, in Figure 12a,b, we present the accuracy and CNOT gate count for $g = 2$ and $g = 3$. Clearly, increasing the value of $g$ not only reduces the error but also increases the number of CNOT gates as seen in Figure 12c,d.

A more complete picture for circuit complexity can be seen from Table 1, where we present the circuit depth, the number of qubits, the number of CNOT gates, the total number of gates, and the number of parameters needed to describe a quantum circuit, using the standard implementation available in Qiskit, for clusters of various sizes. Thus, as per Figure 11, when $g = 1$ is used for these calculations, the resources needed are as dictated by the second row (i.e., $[H_2]_2$) in Table 1. This provides a significant reduction in the resources needed to perform these computations, as noted in the figures mentioned above.

Thus, in general, choice of maximum $g$ and edge length cutoff would be based on a compromise between the desired accuracy and efficiency. However, in all cases, it appears that the current algorithm presents a powerful mode to reduce the quantum circuit depth problem in quantum computing, as summarized by Figure 13.

Table 1. Quantum Circuit Resource Requirements When Standard Techniques Are Used: Illustration for a Family of H$_2$ Clusters

<table>
<thead>
<tr>
<th>system</th>
<th>circuit depth</th>
<th>number of qubits</th>
<th>number of CNOT gates</th>
<th>total number of gates</th>
<th>number of parameters</th>
</tr>
</thead>
<tbody>
<tr>
<td>$H_2$</td>
<td>11</td>
<td>2</td>
<td>3</td>
<td>16</td>
<td>3</td>
</tr>
<tr>
<td>$[H_2]_2$</td>
<td>924</td>
<td>6</td>
<td>615</td>
<td>1217</td>
<td>26</td>
</tr>
<tr>
<td>$[H_2]_3$</td>
<td>5920</td>
<td>10</td>
<td>4684</td>
<td>7370</td>
<td>117</td>
</tr>
<tr>
<td>$[H_2]_4$</td>
<td>21 361</td>
<td>14</td>
<td>16 285</td>
<td>27 021</td>
<td>360</td>
</tr>
<tr>
<td>$[H_2]_5$</td>
<td>57 402</td>
<td>18</td>
<td>47 312</td>
<td>70 204</td>
<td>875</td>
</tr>
<tr>
<td>$[H_2]_6$</td>
<td>128 469</td>
<td>22</td>
<td>107 190</td>
<td>156 081</td>
<td>1818</td>
</tr>
<tr>
<td>$[H_2]_7$</td>
<td>253 846</td>
<td>26</td>
<td>205 192</td>
<td>313 143</td>
<td>3381</td>
</tr>
<tr>
<td>$[H_2]_8$</td>
<td>458 233</td>
<td>30</td>
<td>389 472</td>
<td>550 279</td>
<td>5792</td>
</tr>
</tbody>
</table>

“Clearly, as the system size grows, the circuit gets extremely complex. The Graph-Q|Q approach reduces resource complexity by several orders of magnitude.

Figure 13. Summary of reduction in CNOT gate-based complexity. The red circles are the standard implementation, whereas the blue and green circles arise from the circuit decomposition scheme discussed here.

5. CONCLUSIONS

The promise of solving exponentially complex problems efficiently using quantum computing hardware and associated quantum computing algorithms software is a rapidly evolving research frontier. While we are in the early stages of this quantum revolution, there are a wide set of scientific and technological areas that can benefit from such developments.
However, true progress in such areas can only be achieved by a rigorous study and understanding of the electronic structure and dynamics of complex materials, thus requiring accurate treatment of electron correlation effects in conjunction with a rigorous treatment of quantum nuclear effects.

Many groups have contributed greatly to the development of new quantum circuit-based methodologies to compute the electronic structure in small molecular systems on quantum hardware. Most of these efforts are deeply hindered by the so-called quantum circuit depth problem where the complexity of the quantum circuit, along with the limited fidelity of the quantum gates currently available on state-of-the-art quantum hardware, leads to an enormous increase in error propagation and in stable implementation. This provides a strong upper bound on the accuracy, system size, and levels of basis functions that can be implemented in currently available quantum hardware.

In this publication, we present a novel procedure to reduce the depth of quantum circuits and reduce the extent to which this influences the quality of results in computing post-Hartree–Fock electronic structure energies. The method is based on molecular fragmentation, and specifically, a molecular system is divided into overlapping fragments through a graph-theoretic technique. This then allows the construction of a series of projection operators, that allow some overall model for quantum computing obtained from an approximation to the unitary evolution of the full system, into separate processes, some of which can be treated on quantum hardware and others on classical hardware. Thus, we develop a procedure for electronic structure that can be spawned on to a potentially large ensemble of classical and quantum hardware systems.

We demonstrate this methodology by computing UCCSD energies for a set of \([\text{H}_2\text{O}]\) clusters, with \(n\) ranging from 4 to 128. We implement our methodology using quantum circuits, and when these quantum circuits are processed on a quantum simulator, we obtain energies in agreement with the correct UCCSD energies in the milli-hartree energy range. We also show that our circuit decomposition approach yields up to \(9\) orders of magnitude reduction in the number of CNOT gates and circuit depth for the larger-sized clusters when compared to a standard quantum circuit implementation as available within Qiskit. Future work will involve the implementation of these methods on quantum hardware.

### ASSOCIATED CONTENT

**Supporting Information**

The Supporting Information is available free of charge at https://pubs.acs.org/doi/10.1021/acs.jctc.1c01303.

A few key ideas regarding our graph-theoretic formalism including the generation of nodes, edges, and the definition of higher-order (\(R\)-rank) simplexes that capture \(R = 1\)-body interactions in the molecular system as per eqs 4 and 10 (PDF).

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**Notes**

The authors declare no competing financial interest.

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